

PERFECT MATCHINGS IN
PLANAR CUBIC GRAPHS

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JOINT WITH

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LOVASZ-PLUMMER CONJECTURE, ~1975

$\exists c > 0$ SUCH THAT

EVERY 2-EDGE-CONNECTED CUBIC GRAPH

G HAS $\geq 2^{c|V(G)|}$

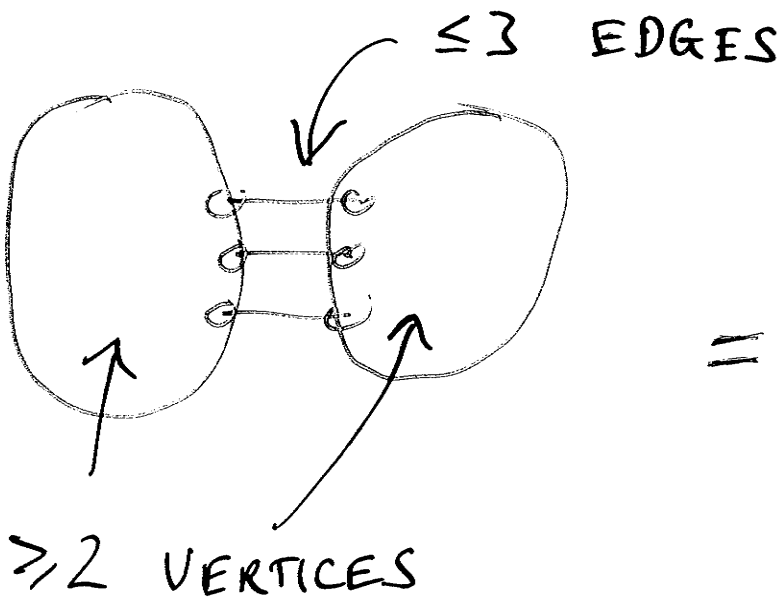
PERFECT MATCHINGS.

THM (LOVASZ-PLUMMER)

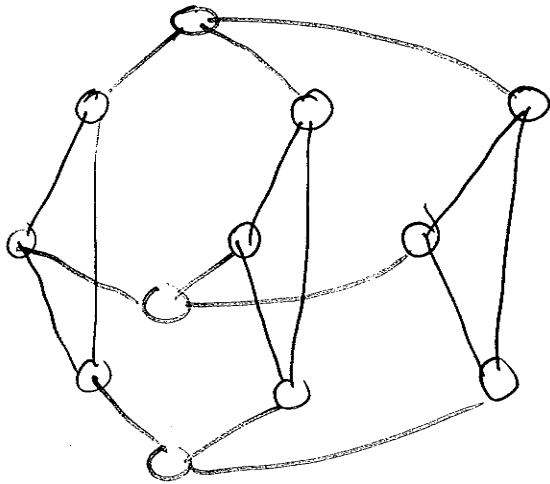
EVERY CYCLICALLY 4-CONNECTED CUBIC

GRAPH G HAS $\geq \frac{1}{2}|V(G)| + 1$

PERFECT MATCHINGS.



= NOT CHC



ONLY 6 PERFECT MATCHINGS

KRAL
2008

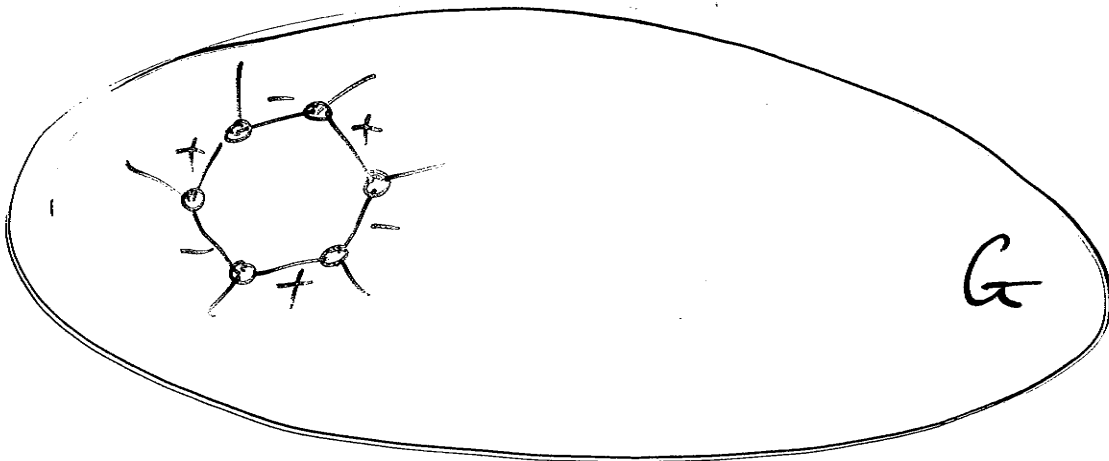
EVERY 2-EDGE-CONNECTED
CUBIC GRAPH G HAS
 $\geq \frac{1}{2} |V(G)|$ PERFECT MATCHINGS

THM EVERY 2-EDGE-CONNECTED PLANAR

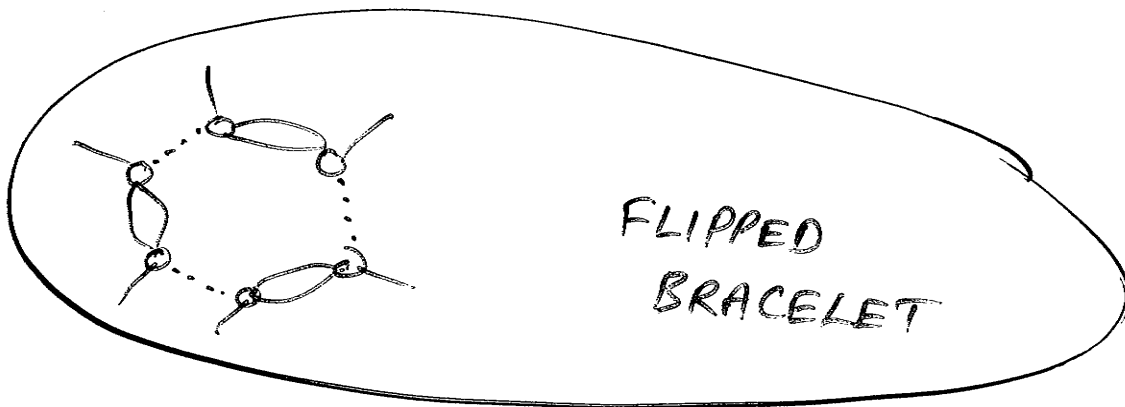
CUBIC GRAPH G HAS

$$\geq 2 \frac{|V(G)|}{655978752}$$

PERFECT MATCHINGS.



BRACELET



LOOK FOR: SET OF DISJOINT BRACELETS
SUCH THAT FLIPPING ANY SUBSET OF THEM
GIVES GRAPH WITH NO CUTEDGE

BLING

$\beta(G)$ = SIZE OF LARGEST SUCH SET

IN PLANAR G ,

PERFECT MATCHINGS $\geq 2^{\beta(G)/3}$

THM IF G IS PLANAR CUBIC, WITH NO
CUTEDGE, THEN

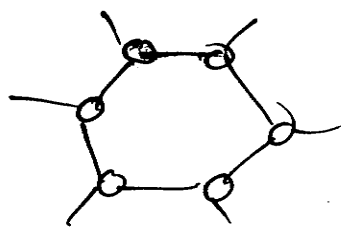
$$\beta(G) \geq \frac{|V(G)|}{218659584}$$

THM: IF G IS CHC PLANAR,

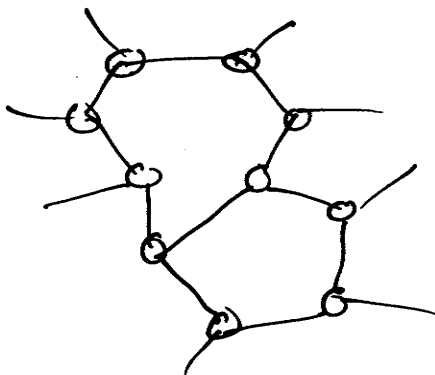
$$\text{THEN } \beta(G) \geq \frac{|V(G)|}{30976}.$$

$$|V(G)| = N. \quad \# \text{ REGIONS} = \frac{N}{2} + 2.$$

DOMINO: EVEN CYCLE BOUNDING ≤ 2 REGION.



OR



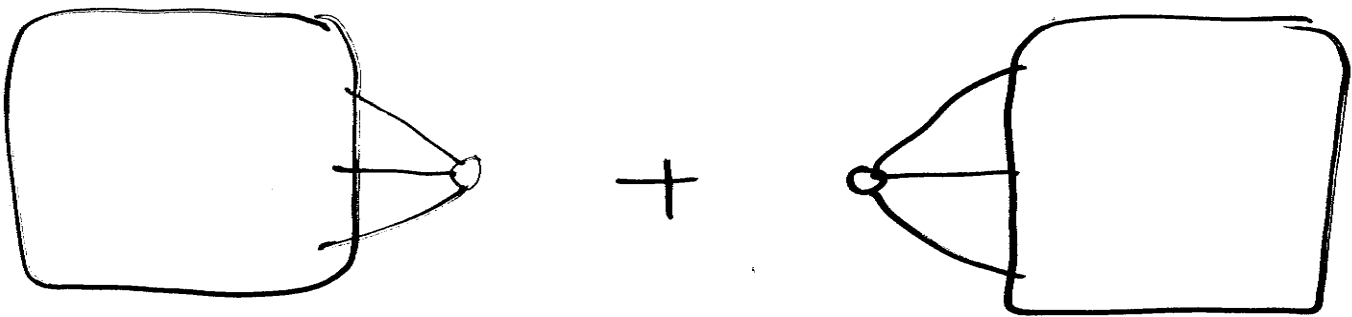
$\frac{N}{4}$ DOMINOES.

$\frac{N}{8}$ ALL OF LENGTH ≤ 15

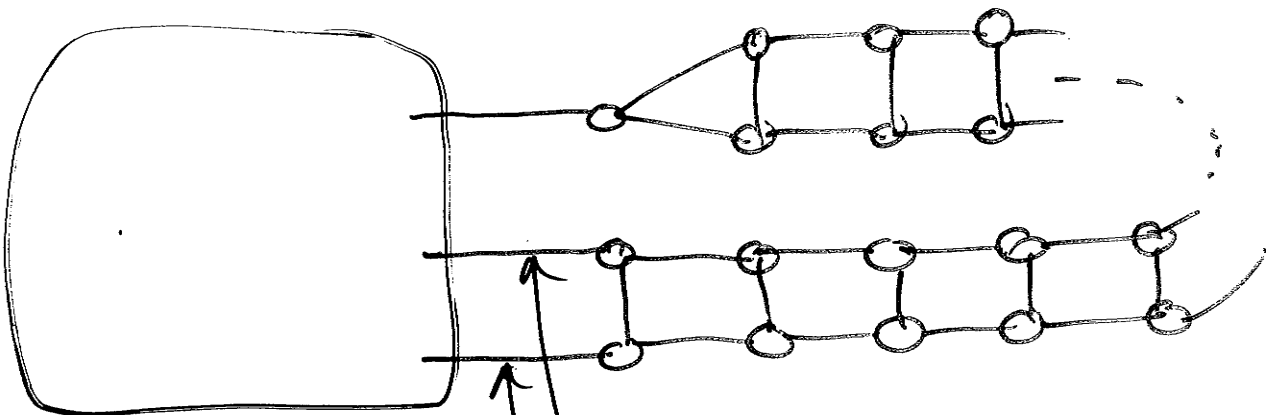
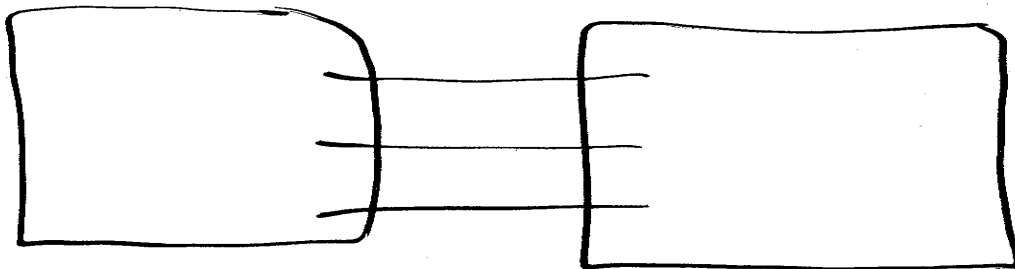
$\frac{N}{32}$ PW VERTEX-DISJOINT

THM IF G IS 3-CONNECTED PLANAR,

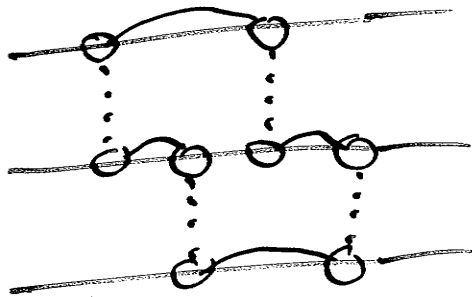
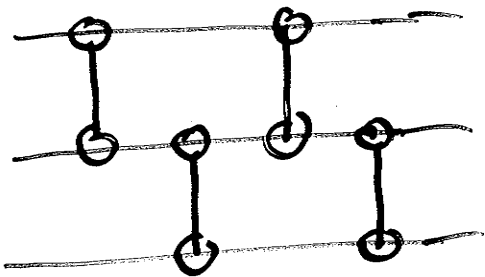
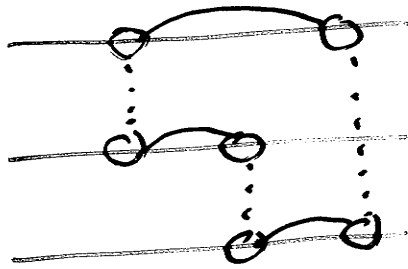
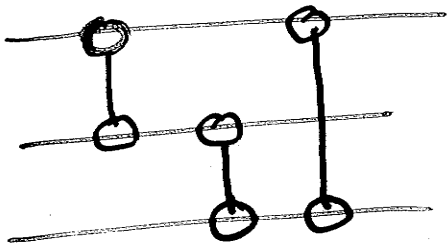
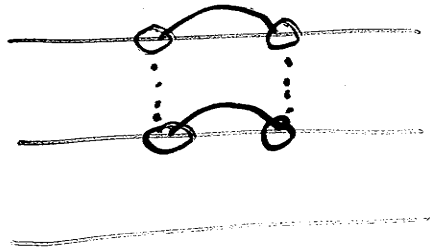
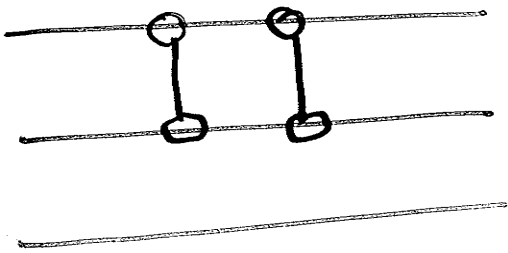
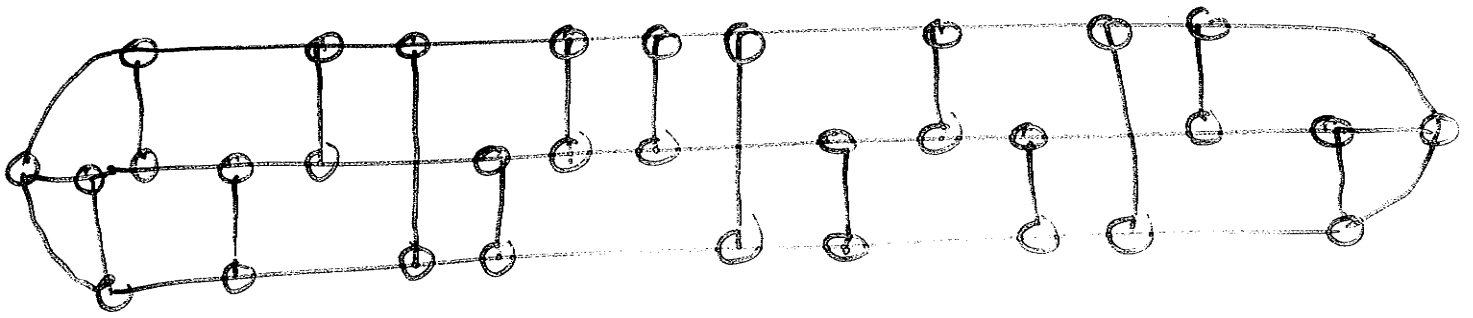
THEN $\beta(G) \geq \frac{|V(G)|}{16819968}$



=



IN UNIQUE PERFECT MATCHINGS

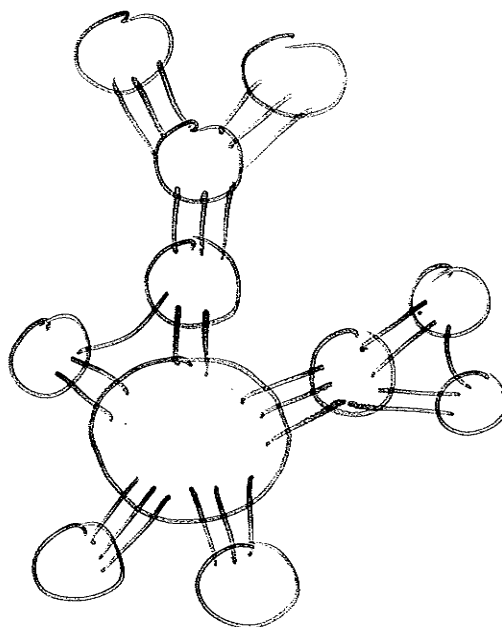
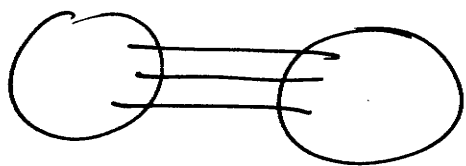
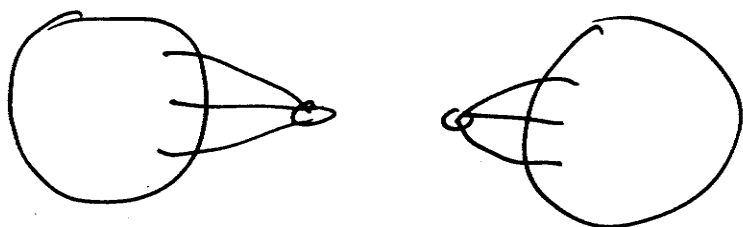


THM (FOWLER, 1998)

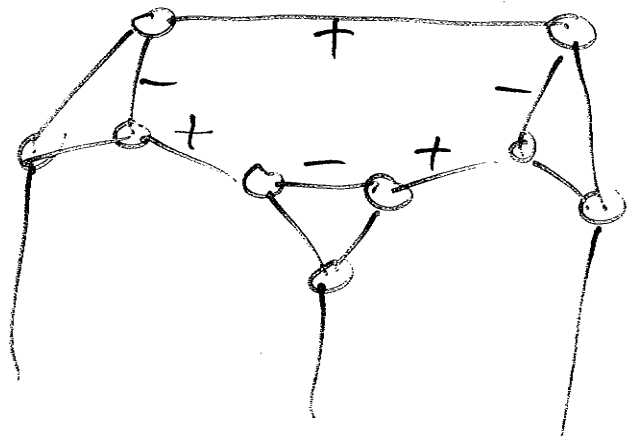
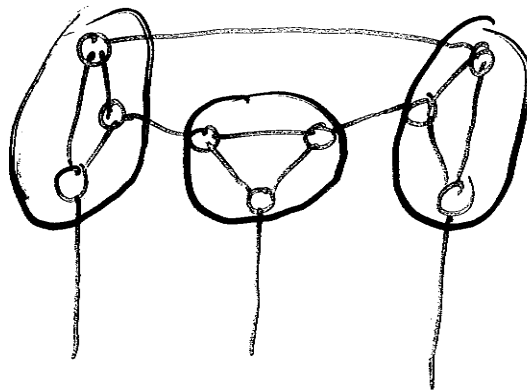
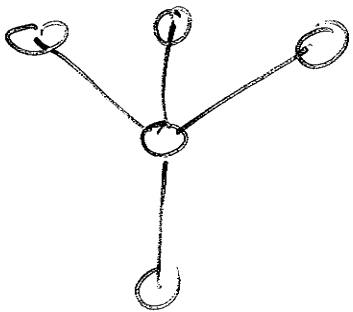
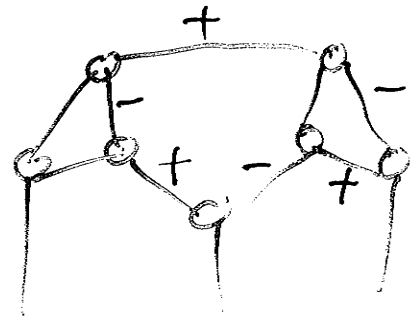
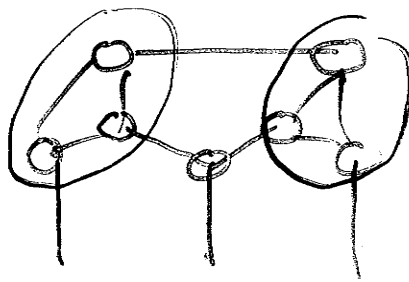
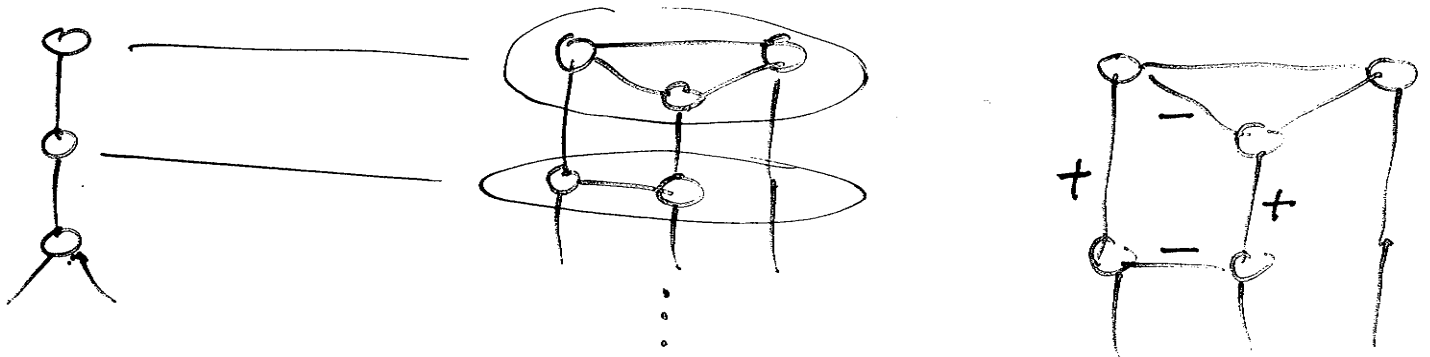
K_4 IS THE ONLY C4C PLANAR

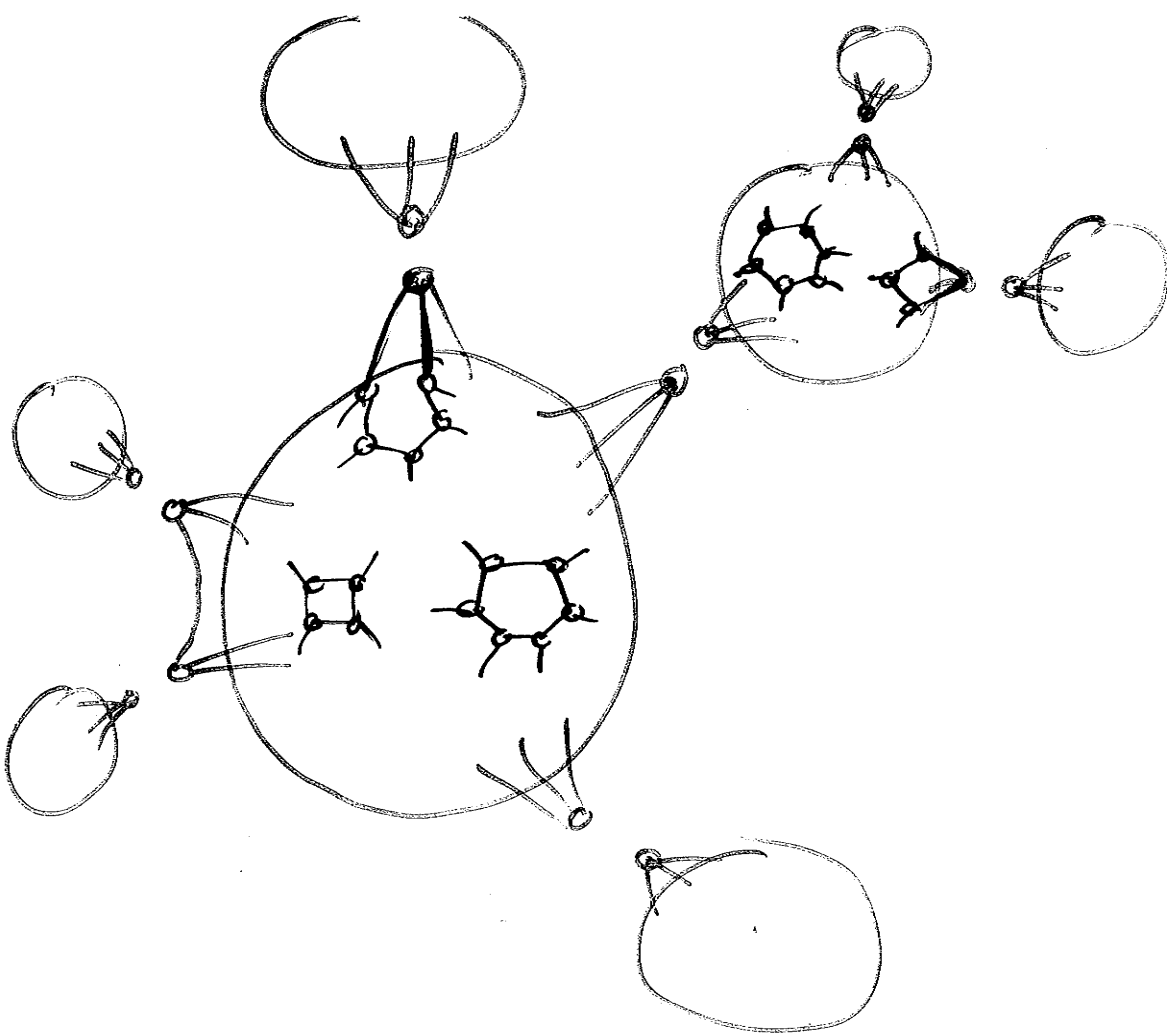
CUBIC GRAPH THAT IS

UNIQUELY 3-EDGE-COLOURABLE.



LEAVES OF FOWLER TREE :





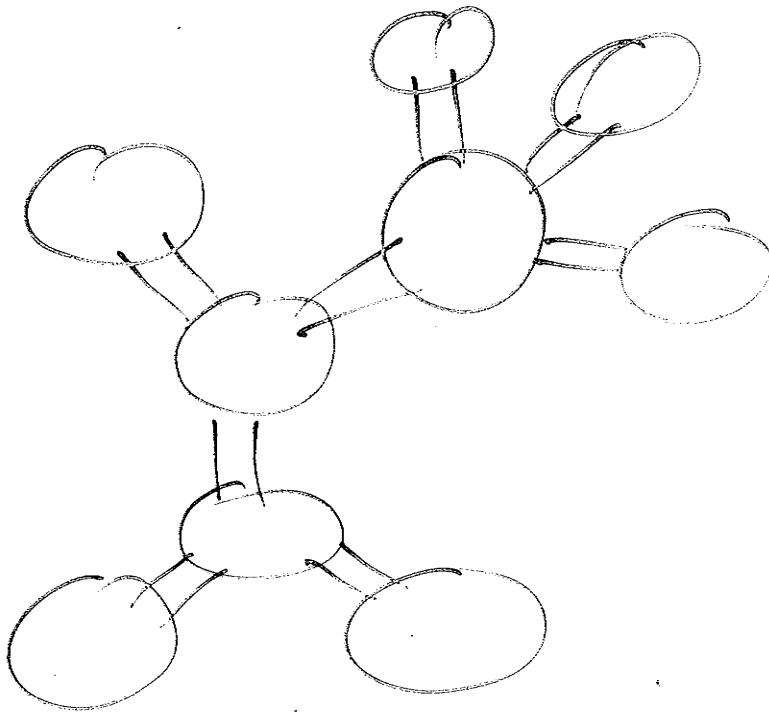
CASE 1 : NO OF CUTS = $O(|V(G)|)$

CASE 2 : $O(|V(G)|)$ PIECES ARE NOT
UNIQUELY 3-COL.

CASE 3 : ALMOST FOWLER ??

REORGANIZE : FIRST, CONTRACT ALL TRIANGLES,
THEN CASES 1, 2, 3

2-CONNECTED CASE



CASE 1: NO OF CUTS

$$O(|V(G)|)$$

CASE 2: $O(|V(G)|)$

CUTS ALL WITH

SAME COLOUR.

